

Sensitivity Analysis

- A few concepts
 - Impulse-Response
- What do we know how to do
 - By MP class
- Related analyses
 - Consistency, Redundancy, & Implied Equalities
- Some foundations
 - Alternative/dual systems
- Some practical considerations
 - Estimating results from partial information

An 'expert' is one who doesn't know more than you but uses slides.

Impulse-Response Queries

	Response	
Impulse	Data	Solution
Data	<i>Drive</i>	<i>Common</i>
Solution	<i>Inverse</i>	<i>Rate of substitution</i>

I never met an optimum I didn't like.

– Milton M. Gutterman

Impulse-Response Queries

	Response	
Impulse	Data	Solution
Data	<i>Drive</i>	<i>Common</i>
Solution	<i>Inverse</i>	<i>Rate of substitution</i>

- At what rate does the objective value change when I perturb some parameter?
For what range is this rate constant (or same functional form)?
- At what rate does the level (or price) change when I perturb some parameter?
For what range is this rate constant (or same functional form)?

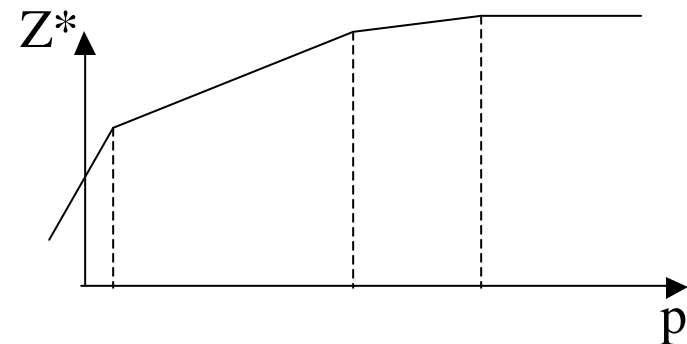
$$\partial^{(\pm)} Z^* / \partial p = k \text{ for } p \in [P-L, P+U]$$

p = parameter

P = current value of p

Z^* = optimal objective value

$[L, U]$ = range of change

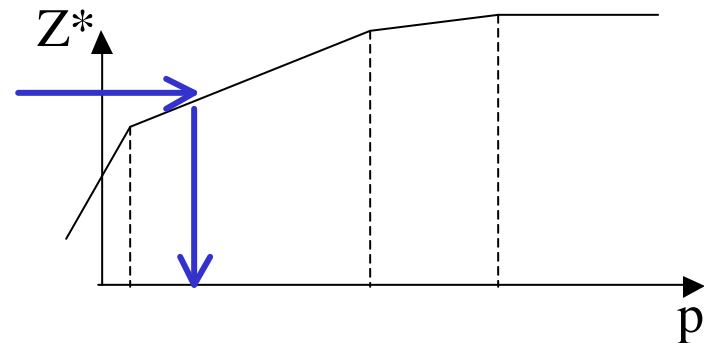


Impulse-Response Queries

	Response	
Impulse	Data	Solution
Data	<i>Drive</i>	<i>Common</i>
Solution	<i>Inverse</i>	<i>Rate of substitution</i>

- How can I change some parameter to cause a 10% decrease in the min cost?
 - e.g., decrease demand or make some inexpensive supply available
- How can I change some parameter to reach specified change in solution?
 - e.g., increase max oxygen to result in more glucose production.

$$\partial^{(\pm)} Z^* / \partial p = k \text{ for } p \in [P-L, P+U]$$

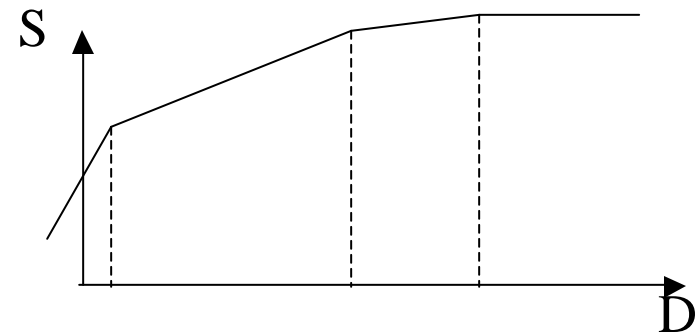


Impulse-Response Queries

	Response	
Impulse	Data	Solution
Data	<i>Drive</i>	<i>Common</i>
Solution	<i>Inverse</i>	<i>Rate of substitution</i>

- How can I change some parameter such that to remain in equilibrium, I must change another (specified) parameter?
 - e.g., decrease demand (D) and increase some (specified) supply (S):

$$\Delta S = k\Delta D$$



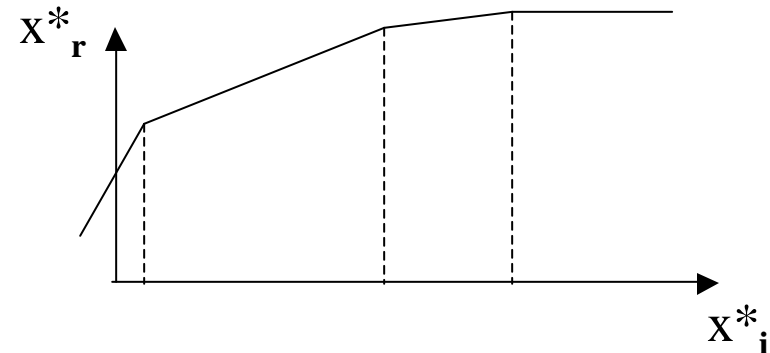
Impulse-Response Queries

	Response	
Impulse	Data	Solution
Data	<i>Drive</i>	<i>Common</i>
Solution	<i>Inverse</i>	<i>Rate of substitution</i>

- How does one solution value change if I force a change in some other?

$$\partial^{(\pm)} x_r^* / \partial x_i^*$$

– Applies to phase-plane analysis.



Simplex Method Uses this Every Iteration

		Nonbasic		
Basic	Level		x_i	
x_r	b_r		a_{ri}	
$-Z$			d_i	

$$x_r = b_r - a_{ri} x_i$$

rate of substitution

Qualitative Analysis

- Given directions of change of parameter, find directions of change of solution
- Find qualitative relations among variables (degrees of separation among metabolites or reactions)
- Find stability properties (not numbers)
- Find pathways of certain interest

Modeling is about insight, not numbers.

– Arthur M. Geoffrion

A Quick Tour of What We Know

- Linear Programming (LP)
- Nonlinear Programming (NLP)
- Integer Programming & Combinatorial Optimization (IP/CO)
- Mixed-Integer Linear Programming (MILP)
- Mixed-Integer Nonlinear Programming (MINLP)

The pure and simple truth is rarely pure and never simple.

– Oscar Wilde

LP

- Basic solution
 - Compatibility theory
- Interior solution
 - Optimal partition
- General case
 - Character of solution

Mostly well understood, but algorithms not perfect

Qualitative analysis strongest for network models, then Leontief

MOLP: Objective space gives important insights

NLP

- Lagrange multipliers
 - Marginal analysis with convexity; “rapid” re-optimization
- Dynamic programming
 - Inherently parametric; needs separability & low dimension
- Pooling problem (bilinear constraints)
 - Exploit geometry to overcome non-convexity
 - Raised new concept – *Essential* objects (pools/reactions)

Sometimes wrong, but never in doubt.

– Michael Evans (Economics forecaster)

IP/CO

- Generally NP-hard
 - Optimizers do not provide automatic support beyond LP
- Special focus on problem structures
 - Scheduling, TSP, covering, packing, ...
- Computational logic
 - Horn clauses: if A then B (single antecedent & consequent)
- New definitions
 - Stability regions; ties with algorithm/heuristic used
- Visualization
 - Diagrammatic; Iconic; Animation

MILP

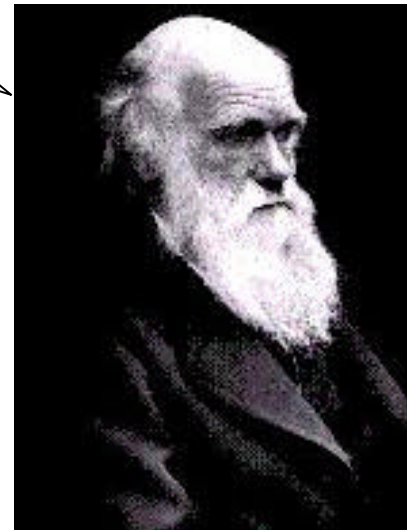
- Loses structural information
 - Preserve logic of IP part (binary variables to control fluxes)

Logical	Algebraic
$x=0 \rightarrow y=0$	$x - y \geq 0$
$x=0 \rightarrow y=1$	$x + y \geq 1$
$x=1 \rightarrow y=0$	$x + y \leq 1$
$x=1 \rightarrow y=1$	$x - y \leq 0$

MINLP

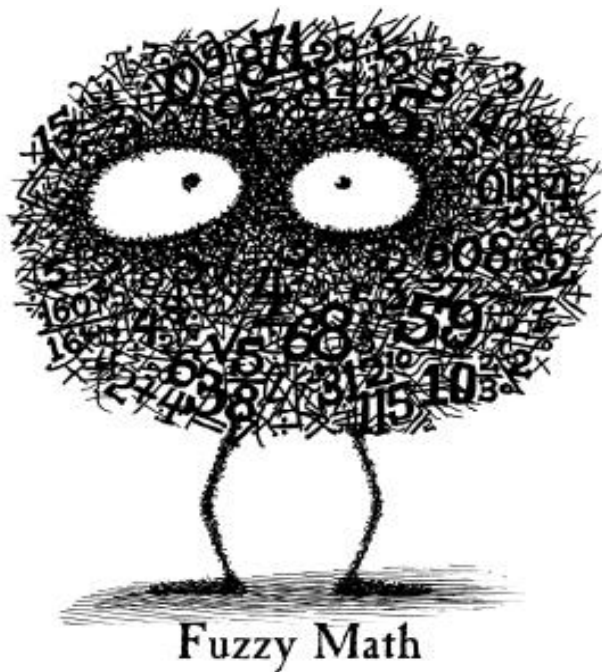
- No theory; few special algorithms
(I. Grossman did some things for specific problems)

*Every new body of discovery
is mathematical in form,
because there is no other
guidance we can have.*



Other Forms

- Multiple objectives
- Goal programs
- Fuzzy programs
- Stochastic programs
- Randomized programs
- Semi-definite programs



Summary of SA Capability

- Linear Programming (LP)
 - Lots known; All queries; Must be careful
- Nonlinear Programming (NLP)
 - Only special cases (convex quadratic; bilinear)
- Integer Programming & Combinatorial Optimization (IP/CO)
 - Hard, but some good results, using logical structure
- Mixed-Integer Linear Programming (MILP)
 - Use IP/CO methods
- Mixed-Integer Nonlinear Programming (MINLP)
 - Uncharted

Consistency, Redundancy, and Implied Equalities

System: $S = \{Ax \geq b\}$

Polyhedron: $P(S) = \{x: Ax \geq b\}$

Subsystems: $S(I) = \{A_i \cdot x \geq b_i \text{ for } i \in I\}$

$S_i = \{A_k \cdot x \geq b_k \text{ for } k \neq i\}$

Inconsistent: $P(S) = \emptyset$

Redundant: $P(S_i) = P(S)$

Strongly Redundant: $x \in P(S_i) \rightarrow A_i \cdot x > b_i$

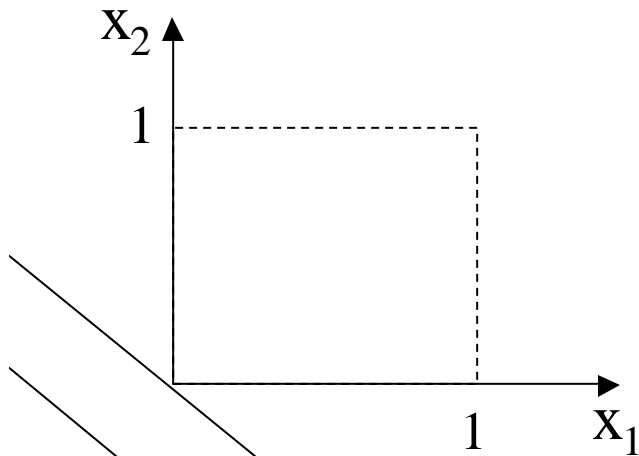
Implied Equality: $A_i \cdot x = b_i \text{ for all } x \in P(S)$

*A model is to an analyst as a magnifying glass is to Sherlock Holmes
– it illuminates clues.*

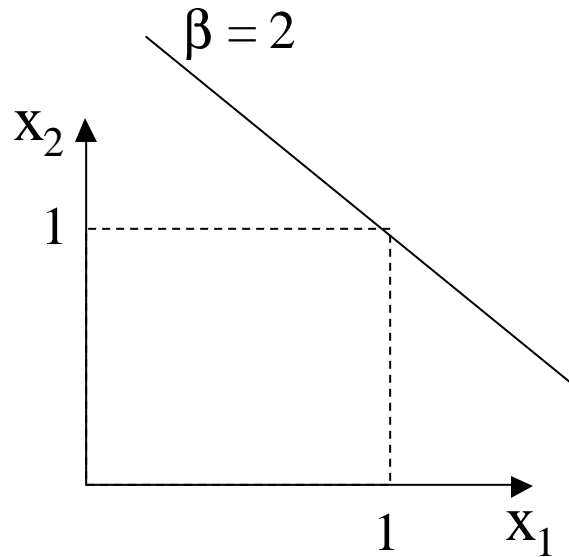
Example

$$S = \{0 \leq x_1, x_2 \leq 1, x_1 + x_2 \geq \beta\}$$

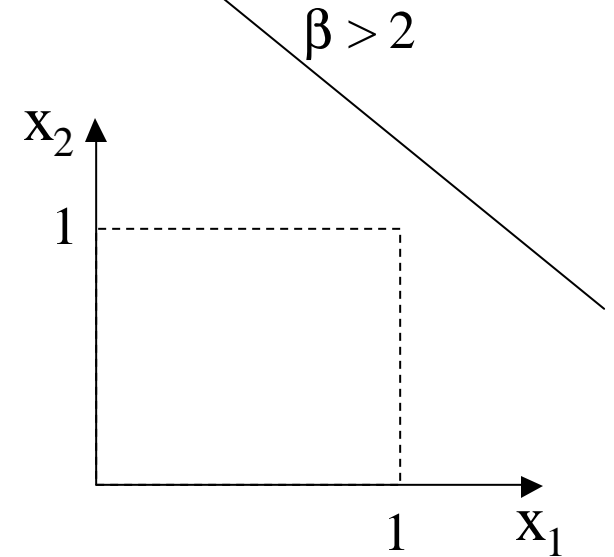
Redundant



Implied equality



Inconsistent



$\beta = 0$
 $\beta > 0$

Foundation = Dual system

$$S^d = \{y \geq 0, yA = 0, yb \geq 0\}$$

Example

$$\begin{array}{rcl} x_1 & & \geq 0 \\ & x_2 & \geq 0 \\ -x_1 & & \geq -1 \\ & -x_2 & \geq -1 \\ x_1 + x_2 & & \geq \beta \end{array} \qquad \begin{array}{l} y_1, y_2, y_3, y_4, y_5 \geq 0 \\ y_1 - y_3 + y_5 = 0 \\ y_2 - y_4 + y_5 = 0 \\ -y_3 - y_4 + \beta y_5 \geq 0 \end{array}$$

A study of economics usually reveals that the best time to buy anything is last year.

– Marty Allen

Certification

Property of S is true $\leftrightarrow S^*$ is consistent

Property of S	S^*
Redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0\}$
Strongly redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0, y_b > 0\}$
Implied equality	$S^d \ \& \ \{y_i > 0, y_b = 0\}$
Inconsistent	$S^d \ \& \ \{y_b > 0\}$

$$y_1, y_2, y_3, y_4, y_5 \geq 0$$

$$y_1 - y_3 + y_5 = 0$$

$$y_2 - y_4 + y_5 = 0$$

$$-y_3 - y_4 + \beta y_5 \geq 0$$

Certification of Redundancy

Property of S is true $\leftrightarrow S^*$ is consistent

Property of S	S^*
Redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0\}$
Strongly redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0, yb > 0\}$
Implied equality	$S^d \ \& \ \{y_i > 0, yb = 0\}$
Inconsistent	$S^d \ \& \ \{yb > 0\}$

$$\beta = 0$$

choose $y = (1, 1, 0, 0, -1)$

$$y_1, y_2, y_3, y_4, \cancel{y_5} \geq 0 \quad y_5 < 0$$

$$y_1 - y_3 + y_5 = 0$$

$$y_2 - y_4 + y_5 = 0$$

$$-y_3 - y_4 \geq 0 \quad \rightarrow y_3 = y_4 = 0$$

Certification of Strong Redundancy

Property of S is true $\leftrightarrow S^*$ is consistent

Property of S	S^*
Redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0\}$
Strongly redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0, yb > 0\}$
Implied equality	$S^d \ \& \ \{y_i > 0, yb = 0\}$
Inconsistent	$S^d \ \& \ \{yb > 0\}$

$$\beta < 0$$

choose $y = (1, 1, 0, 0, -1)$

certifies strong redundancy
because $yb = -\beta > 0$

$$y_1, y_2, y_3, y_4, \cancel{y_5} \geq 0 \quad y_5 < 0$$

$$y_1 - y_3 + y_5 = 0$$

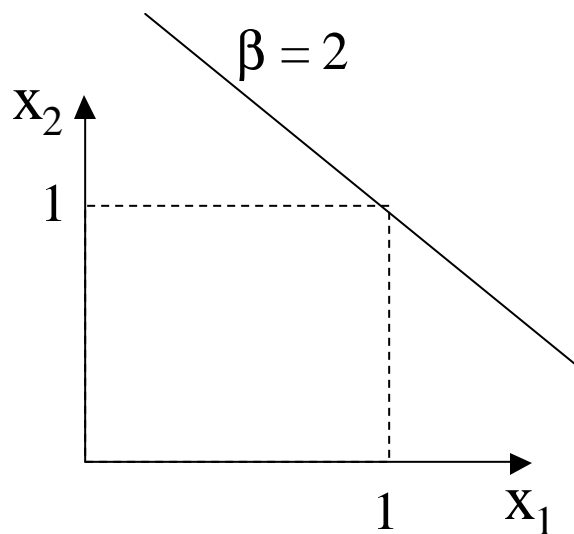
$$y_2 - y_4 + y_5 = 0$$

$$-y_3 - y_4 + \beta y_5 \geq 0$$

Certification of Implied Equality

Property of S is true $\leftrightarrow S^*$ is consistent

Property of S	S^*
Redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0\}$
Strongly redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0, y_b > 0\}$
Implied equality	$S^d \ \& \ \{y_i > 0, y_b = 0\}$
Inconsistent	$S^d \ \& \ \{y_b > 0\}$



$$y_1, y_2, y_3, y_4, y_5 \geq 0$$

$$y_1 - y_3 + y_5 = 0$$

$$y_2 - y_4 + y_5 = 0$$

$$-y_3 - y_4 + \beta y_5 \geq 0$$

Certification of Implied Equality

Property of S is true $\leftrightarrow S^*$ is consistent

Property of S	S^*
Redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0\}$
Strongly redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0, y_b > 0\}$
Implied equality	$S^d \ \& \ \{y_i > 0, y_b = 0\}$
Inconsistent	$S^d \ \& \ \{y_b > 0\}$

$$\beta = 2$$

choose $y = (0, 0, 1, 1, 1)$

plays no role in implied equality

$$y_1, y_2, y_3, y_4, y_5 \geq 0$$

$$y_1 - y_3 + y_5 = 0$$

$$y_2 - y_4 + y_5 = 0$$

$$-y_3 - y_4 + 2y_5 \geq 0$$

Certification of Inconsistency

Property of S is true $\leftrightarrow S^*$ is consistent

Property of S	S^*
Redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0\}$
Strongly redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0, y_b > 0\}$
Implied equality	$S^d \ \& \ \{y_i > 0, y_b = 0\}$
Inconsistent	$S^d \ \& \ \{y_b > 0\}$

$$\beta > 2$$

choose $y = (0, 0, 1, 1, 1)$

plays no role in inconsistency

$$y_1, y_2, y_3, y_4, y_5 \geq 0$$

$$y_1 - y_3 + y_5 = 0$$

$$y_2 - y_4 + y_5 = 0$$

$$-y_3 - y_4 + \beta y_5 \geq 0$$

Certification of Inconsistency

Property of S is true $\leftrightarrow S^*$ is consistent

Property of S	S^*
Redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0\}$
Strongly redundant	$S^d \setminus \{y_i \geq 0\} \ \& \ \{y_i < 0, y_b > 0\}$
Implied equality	$S^d \ \& \ \{y_i > 0, y_b = 0\}$
Inconsistent	$S^d \ \& \ \{y_b > 0\}$

$$\beta > 2$$

choose $y = (0, 0, 1, 1, 1)$

certifies inconsistency

because $y_b = \beta - 2 > 0$

$$y_1, y_2, y_3, y_4, y_5 \geq 0$$

$$y_1 - y_3 + y_5 = 0$$

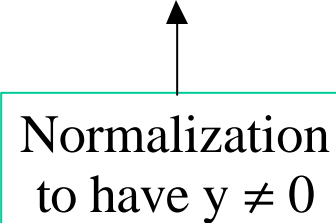
$$y_2 - y_4 + y_5 = 0$$

$$-y_3 - y_4 + \beta y_5 \geq 0$$

Certificates Obtained by LP

$$\max yb: y \in P(S^{\sim}), \quad \sum_i y_i = 1$$

Normalization
to have $y \neq 0$



Interior Solutions Certify All at Once with $s(y)$

$$y^* \in \operatorname{argmax}\{yb: yA = 0, y \geq 0, \sum_i y_i = 1\}$$

\Rightarrow every i for which $y^*_i > 0$ is an implied equality

y^* interior $\Rightarrow \sigma(y) =$ set of *all* implied equalities of $Ax \geq b$